

### Mohammad Farshi

Department of Computer Science, Yazd University

1394-2



Randomized Algorithms

Course Outline

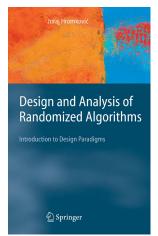
Textbook

Prerequisties

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### Textbook:

Juraj Hromkovic-Design and Analysis of Randomized Algorithms- Introduction to Design Paradigms-Springer (2005)





Randomized Algorithms

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What is Bandomized

### Prerequisites:



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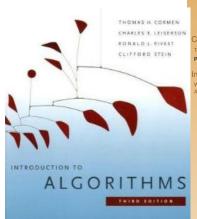
What is Randomized Algorithms about?

### What you need to know:

- Basic Algorithms and Algorithm Analysis: O, Θ notations, sorting, searching.
- Basic Data Structures: Priority Queue (Heap), Binary Search Tree, ... and their analysis.
- Basic Probability theory: Expected value, ...

# Prerequisites:







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Introduction



### Randomized Algorithms

### Applying randomness to:

- designing algorithms that are very fast is most cases, but slow in very rare cases (in comparison to deterministic algorithms).
- designing algorithms that are very fast, but in most cases the output of the algorithm is true, but in some rare cases the output is false.
- prove deterministic results (existence of something
- get rid of the adversary, or make the algorithm independent to the shape of input.

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### Randomized Algorithms

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### **Applications**

- designing practical algorithms (very simple to implement, very fast in practice),
- make complicated problems very easy,
- an step forward to a deterministic algorithm for a problem,
- proving the existence of something (if the probability of an event is non-zero, it can happens),
- and so on.

### An Example:

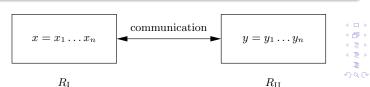
- We have two computers  $R_I$  and  $R_{II}$  that are very far apart.
- At the beginning both have a database with the same content (say,  $n=10^{16}$  bits).
- The contents of these databases dynamically developed simultaneously in both databases.
- After some time, we want to check whether  $R_I$  and  $R_{II}$  contain the same data.
- Goal: Design a communication protocol between R<sub>I</sub> and R<sub>II</sub> to check this.



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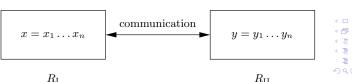


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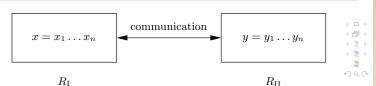
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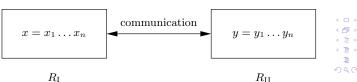


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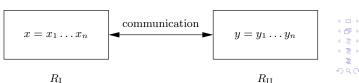


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### The solution:

- The complexity of the communication protocol = the number of bits exchange between computers.
- There exist no deterministic protocol that solves this task by communicating < n bits.</li>
- Sending  $n = 10^{16}$  bits safely is a practically nontrivial task, so one would probably not do it in this way.
- Randomness can help here!

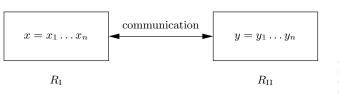
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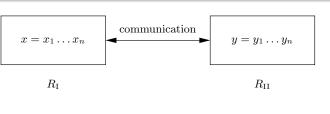
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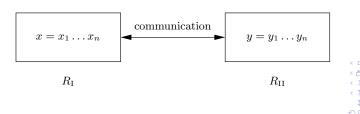


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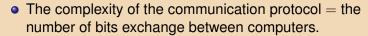
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communication
x = x_1 \dots x_n
                                                              y = y_1 \dots y_n
        R_{\rm T}
                                                                     R_{\rm II}
```

### **Randomized Protocol for Equality**

- Initial situation:  $R_I$  has a sequence x of n bits,  $x=x_1\cdots x_n$ , and  $R_{II}$  has a sequence y of n bits  $y=y_1\cdots y_n$ .
- Phase 1:  $R_I$  chooses uniformly a prime p from the interval  $[2, n^2]$  at random.
- Phase 2:  $R_I$  computes the integer  $s = Number(x) \mod p$  and sends the binary representations of s and p to  $R_{II}$ . Observe that  $s \le p < n^2$  and so each of these integers can be represented by  $\lceil \log_2 n^2 \rceil$  bits.
- Phase 3: After reading s and p,  $R_{II}$  computes the number  $q = Number(y) \mod p$ . If  $q \neq s$ , then  $R_{II}$  outputs  $x \neq y$ . If q = s, then  $R_{II}$  outputs x = y.



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### **Randomized Protocol for Equality**

- Complexity of the protocol: Since  $s \le p < n^2$ , the length of the message is  $\le 2\lceil \log_2 n^2 \rceil \le 4\log_2 n$ . For  $n=10^{16}$ , it is 256.
- Reliability (error probability) of the protocol



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- Reliability (error probability) of the protocol: The answer of the protocol is not always correct: For instance, if x=01111 and y=10110, i.e., Number(x)=15 and Number(y)=22, then the choice of the prime 7 from the set  $\{2,3,5,7,11,13,17,19,23\}$  yields the wrong answer, because  $15 \mod 7 = 1 = 22 \mod 7$ .



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$$PRIM(n^2) = \{p \text{ is a prime}|p \leq n^2\}$$



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 $\begin{array}{c|c} \text{bad} & \text{good primes for} \\ \text{primes} & \text{the input } (x,y) \\ \\ \text{for } (x,y) & \text{all primes} \leq n^2 \end{array}$ 

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Error probability for  $(x,y)=\frac{\# \text{ bad primes for } (x,y)}{Prim(n^2)}.$ 



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If we show that # bad primes  $\leq n-1$ , then the error probability

$$\leq \frac{n-1}{Prim(n^2)} \leq \frac{n-1}{n^2/\log n^2} \leq \frac{\log n^2}{n}.$$

For  $n = 10^{16}$ , this probability is  $0.36892 \times 10^{-14}$ 



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Randomized Protocol for Equality: # bad primes

# # bad primes

• p is a bad prime if  $x \neq y$  but

$$Number(x) \bmod p = Number(y) \bmod p.$$

i.e. p divides w = |Number(x) - Number(y)|.

- Let  $w = p_1^{i_1} \times p_2^{i_2} \times \cdots \times p_k^{i_k}$  is prime factorization of w  $(p_i < p_{i+1})$ .
- p<sub>i</sub>s are all bad primes.
- claim:  $k \le n-1$ . Assume  $k \ge n$ .

$$w = p_1^{i_1} \times p_2^{i_2} \times \dots \times p_k^{i_k}$$

$$\geq p_1 \times p_2 \times \dots \times p_k$$

$$> 1 \times 2 \times \dots \times n = n! > 2^n(Contradiction!)$$



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### **Randomized Protocol for Equality**

- Reliability (error probability) of the protocol: not satisfying?
- Run the protocol 10 times, the error probability  $\leq$

$$\leq \left(\frac{n-1}{Prim(n^2)}\right)^{10} \leq \left(\frac{n-1}{n^2/\log n^2}\right)^{10} \leq \frac{2^{10}\log n^{10}}{n^{10}}.$$

For  $n = 10^{16}$ , this probability is  $0.4717 \times 10^{-141}$ 



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